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DL-Lite is a description logic and a subfragment of OWL-DL;

- **DL-Lite is less expressive than OWL-DL;**
- DL-Lite does not allow qualified role number restrictions, nor nominals, nor role hierarchies.

Several reasoning problems have been shown tractable when DL-Lite is involved; unfortunately tractability only applies when there is no restrictions on the size of models (infinite models are possible);

A DL-Lite knowledge base consists of a set

of axioms *C* is subsumed in *D*,

where C are (complex) concepts formed:

- from a set of Concept Names;
- through the complentation of a concept;
- through the conjunction of a pair of concepts;
- through the union of a pair of concepts;

 through the restiction of a role or its inverse (binary relation) by a base number of successors;

A DL-Lite horn knowledge base consists of a set of axioms *C* is subsumed in *D*,

where C, D are (complex) concepts formed:

- from a set of Concept Names;
- through the conjunction of a pair of concepts;

 through the restiction of a role or its inverse (binary relation) by a base number of successors;

from the nill concept (only occurring as D);

Deciding Satisfiability of knowledge bases expressed in general DL-Lite (or DL-Lite bool) is in general a NP-Time problem when sentences are interpreted on unrestricted structures;

The same problem of the horn fragment of DL-Lite, is however tractable for this class of structures;

The exact upper bound of the complexity for deciding satisfiability of general DL-Lite when only finite models are involved is not known, although the problem is known to be NP-hard. We determine the unknown bound here.

Frequency Constraints:

Computing the frequency of Itemsets in a database is a common task in mining frequent patterns;

- An boolean item in a transaction can be easily regarded as a a conceptual content of an object in some world;
- A frequency contraint is a pattern expressing bounds on the frequency of occurrence of a set of items in a database;

Example: *freq* (A, B): [9/10, 1]

expresses that itemset A,B occurs at a rate of 90% of the transactions in the database.

Deciding the possibility of existence of a database fullfilling a set of frequency constraints is known to be an NP-complete problem (Calders, 2002).

First question:

There is some correspondence among DL-Lite and frequency constraints satisfiability problems?

We provide here a strong positive answer.

Second question:

How hard results the integration of DL-Lite knowledge bases with frequency constraints? We answer this question here.

Integration Problem:

Input: A DL-Lite knowledge base; A set of frequency constraints; A DL-Lite concept;

Output: There exists at least one finite model where concept C is non-empty and the axioms and constraints are satisfied?

Integration Problem:

Relevance of Use:

First Setting: Pure Ontological Description;

Example:

Defining the concept Married People.

Any extension of this concept must consists of an even number of members.

<u>Conclusion</u>: Concepts designing unordered structured objects can be defined using frequency constraints.

Integration Problem:

Relevance of Use:

Second Setting: Contextualized Reasoning;

Example:

Deseases as Concepts in a Medical Knowledge Base.

Problem of Reasoning with a real context of frequencies of occurrence of diseases given.

Integration Problem:

Relevance of Use:

Second Setting: Summarizing and Semantic Indexing;

Example:

Summaries may be based on frequencies of occurrences of complex conceptual descriptions. Semantic Indexing often requires the knowledge of the frequency of occurrence of concepts defined in an ontology and reasoning on them.

First Result:

There is a linear time translation of any knowledge base expressed in DL-Lite to a set of frequency constraints, which preserves the finite satisfiability of concepts.

Consequence: DL-Lite concept satisfiability with respect to finite models in NP.

CRUX: Flattening Finite Models

Principles:

Any concept name C in an axiom is a flat item;

Any number restriction in an axion is an item to flatten;

Any conjunction of many concepts in an axiom is an itemset;

Any complex concept is subject of flattening;

Any subsumption axiom C is subsumed in D is translated as freq (C, \overline{D}) = 0;



Since any union of a pair of concepts can always be expressed through a negation of the negation of both concepts to unite, a union can always be flattened. Flattening Complex Concepts: Union



Flattening Number Restrictions: Simplest Case

Role Reification; Substitution of pairs with Counting items; Two Domains as

items: *mon* and *bin*



Theorem 1:

The translation described before receives a knowledge base expressed in DL-Lite and a DL-Lite concept and produces a set of frequency constraints, which preserves the satisfiability of the concept received.

Second Result:

There is a linear time translation of any set of frequency constraints to a knowledge base and a concept expressed in the horn fragment of DL-Lite, which preserves finite satisfiability.

Consequence: DL-Lite *horn* concept satisfiability with respect to finite models is an NP-hard problem (already known).



Itemset: {i1, i2, i3} as a DL-Lite Concept

Constraint: *freq* (Cj, D): [2/5, 3/5]



We have: freq (Cj, D) = 1 / 2 A Frequency Axiom As A DL-Lite Structure



A Frequency Axiom As A DL-Lite Structure

 $freq (Cj, D) \le 3 / 5$ Sj = Im(D, 1...>3) superset Tj = Im(Cj, 1-->5)

Frequency Upper Bound

Constraint: freq (Cj, D): [2/5, 3/5]



Theorem 2:

The translation described before receives a a set of frequency constraints and produces a knowledge base expressed in DL-Lite *horn* and a DL-Lite *horn* concept, such that if there exists a database satisfying the constraints then the concept is finitely satisfiable.

Conclusions:

• The complexity bounds for the concept satisfiability problem of DL-Lite *horn* and *bool logics* is the same. Both problems are NP-complete.

•The enhancing of DL-Lite knowledge bases with frequency constraints does not increase the complexity of the concept satisfiability problem. Two methods emerge as alternatives.

Insights on Tractability:

Since any DL-Lite knowledge base can be safely translated into a set of frequency constraints in linear time, the notion of *closedness* (Pasquier et al. 2005) may be a interesting candidate for a *new semantics* for DL-Lite. With the assumption of closedness for itemsets, implication results tractable.

Future Work:

• Exploring the complexity with the addition of role hierarchies and nominals (the problem for the case of qualified number restriction is already known as EXPTIME cpomplete).

•The first translation process (the flattening process) constitutes the base of a method of adornment of graph databases for mining possible explaining ontologies from them.

Thank You All